Exercise 1: Communication forms in distributed systems

There are basically two types of communication in distributed systems: reliable message streams and unreliable datagrams. What are the main characteristics of both forms of communication and what are their advantages and disadvantages? Name possible areas of application.

Exercise 2: Conversion of data formats

When communicating in distributed systems, it is generally necessary to convert data structures stored internally in the processes into a format suitable for transmission. For what reasons is this indispensable? Discuss necessary measures for different data types.

Exercise 3: Semantics of client/server communication

a) Client/server communication can cause problems if messages are lost during transmission between client and server. The handling of such errors can be different, resulting in different semantics for communication:

- **at least once** - the operation is executed at least once,
- **at most once** - the operation is executed at most once, and
- **exactly once** - the operation is executed exactly once.

The following figure shows state diagrams for client and server with an implementation of **at least once** semantics:

![State Diagram](image)

Sketch analog state diagrams for implementations of the **at most once** and **exactly once** semantics using incoming messages and timeout events as above. Use sequence numbers to identify individual requests for the **exactly once** semantics.

b) For the following applications, consider whether **at least once** or **at most once** semantics is more appropriate:

- pressing an elevator button,
- translating a program,
Exercise 4: Request/reply protocol

For a typical client-server protocol in request/reply style, it makes sense to use a protocol based on datagrams. Why? Design such a protocol.

Pay particular attention to all conceivable error scenarios and how they should be handled. The protocol should be as lightweight as possible. For example, the datagram layer may offer the following operations as a service:

- `send(address, data)` - asynchronous, unreliable transmission of data,
- `recv(address, data, timeout)` - blocking receive of a packet with timeout.

The following operations are to be implemented:

- `doOperation(address, request, reply)` - synchronous sending of a request to the server,
- `getRequest(address, request)` - blocking reception of a request (server),
- `sendReply(address, reply)` - sending the reply (server)

You may assume that fragmentation of the messages to be transmitted is not necessary. Furthermore, it is guaranteed that no data packets are corrupted. However, any packet loss or non-compliance with the packet sequence must be taken into account.

Exercise 5: Middleware

A reliable multicast service allows a sender to deliver reliable messages to multiple recipients. Is such a service part of a middleware layer or should it be part of an underlying layer?

Exercise 6: Transparency of RPC

If Remote Procedure Call (RPC) is used, it does not matter whether the called procedure (i.e., the server process) is located on the local computer node or on another node. But what happens when the procedure executes a system call? What problems could this cause, and how could they be addressed?